

Mathematical Logics

Propositional Logic *

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Reference(s):

- Francesco Berto,
Logica da zero a
Gödel, Laterza, 2018
(capitolo 1)

**Originally by Luciano Serafini and Chiara Ghidini
Modified by Fausto Giunchiglia and Mattia Fumagalli*

1. Intuition

2. Language

3. Satisfiability

4. Validity and unsatisfiability

5. Logical consequence and equivalence

6. Axioms and theories

Definition (Propositional alphabet)

Logical symbols $\neg, \wedge, \vee, \supset, \equiv$

Non logical symbols A set \mathbf{PROP} of symbols P called **propositional variables**

Separator symbols “(” and “)”

Definition (Well formed formulas (or simply formulas))

- every $P \in \mathbf{PROP}$ is an **atomic formula**
- every atomic formula is a **formula**
- if A and B are formulas then $\neg A$, $A \wedge B$, $A \vee B$, $A \supset B$, and $A \equiv B$ are **formulas**

Propositional variables and constants

- Propositional **constants**:
 - “B. Obama is drinking a bier”
 - “The U.S.A. president is drinking a bier”,
 - “B. Obama si sta facendo una birra”
- Propositional **variables**: A, B, C, AA, ..., where a propositional variable can be substituted with a propositional constant, a propositional variable or in general any well formed (propositional) formula

Example ((non) formulas)

Formulas	Non formulas
$P \rightarrow Q$	$P Q$
$P \rightarrow (Q \rightarrow R)$	$(P \rightarrow \wedge ((Q \rightarrow R)$
$P \wedge Q \rightarrow R$	$P \wedge Q \rightarrow \neg R \neg$

Problem

How do we read the formula $P \wedge Q \rightarrow R$?

The formula $P \wedge Q \rightarrow R$ can be read in two ways:

- 1 $(P \wedge Q) \rightarrow R$
- 2 $P \wedge (Q \rightarrow R)$

Symbol priority

\neg has higher priority, then $\wedge, \vee, \rightarrow, \equiv$ with decreasing priority

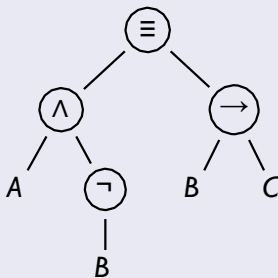
Parenthesis can be used to stress or change the priority.

Symbol	Priority
\neg	1
\wedge	2
\vee	3
\rightarrow	4
\equiv	5

Formulas as trees

A formula can be seen as a tree. Leaf nodes are associated to propositional variables, while intermediate (non-leaf) nodes are associated to connectives.

For instance the formula $(A \wedge \neg B) \equiv (B \rightarrow C)$ can be represented as the tree



Definition

(Proper) Subformula

- A is a **subformula** of itself
- A and B are **subformulas** of $A \wedge B$, $A \vee B$, $A \supset B$, $eA \equiv B$
- A is a subformula of $\neg A$
- if A is a subformula of B and B is a subformula of C , then A is a subformula of C .
- A is a **proper subformula** of B if A is a subformula of B and A is different from B .

Remark

The subformulas of a formula represented as a tree correspond to all the different subtrees of the tree associated to the formula, one for each node.

Example

The subformulas of $(p \rightarrow (q \vee r)) \rightarrow (p \wedge \neg p)$ are

$(p \rightarrow (q \vee r)) \rightarrow (p \wedge \neg p)$

$(p \rightarrow (q \vee r))$

$p \wedge \neg p$

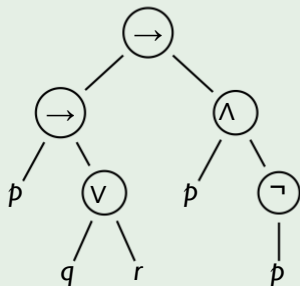
p

$\neg p$

$q \vee r$

q

r



Proposition

Every formula has a finite number of subformulas

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