# Mathematical Logics Modal Logic: K and more\*

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Formulas can be used to shape the "form" of the structure, namely to impose properties on the accessibility relation R.

### Examples

- **Temporal logic**: if the accessibility relation is supposed to represent a temporal relation, and wRw<sup>1</sup> means that w<sup>1</sup> is a future world w.r.t. w, then R must be a transitive relation. That is if w<sup>1</sup> is a future world of w, then any future world of w<sup>1</sup> is also a future world of w.
- Logic of knowledge: if the accessibility relation is used to represent the knowledge of an agent A, and wRw<sup>1</sup> represents the fact that w<sup>1</sup> is a possible situation coherent with its actual situation w, then R must be reflexive, since w is always coherent with itself.

## Typical Properties of R

The following table summarizes the most relevant properties of the accessibility relation, which have been studied in modal logic, and for which it has been provided a sound and complete axiomatization

### Properties of R

R is reflexive R is transitive R is symmetric R is Euclidean R is serial R is weakly dense R is partly functional R is functional R is weakly connected R is weakly directed

```
\forall w.R(w,w)
\forall w \ v \ u.(R \ (w, v) \land R \ (v, u) \supset R \ (w, u))
\forall w \ v \ .(R \ (w, v \ ) \supset R \ (v, w \ ))
\forall w \ v \ u.(R \ (w, v) \land R \ (w, u) \supset R \ (v, u))
\forall w. \exists v R (w, v)
\forall w v . R (w, v) \supset \exists u . (R (w, u) \land R (u, v))
\forall w \ v \ u.(R \ (w, v)) \land R \ (v, u) \supset v = u)
\forall w \exists ! v. R (w, v)
\forall u \ v \ w.(R \ (u, v)) \land R \ (u, w)) \supset
             R(v, w) \lor v = w \lor R(w, v))
\forall u \ v \ w.(R \ (u, v)) \land R \ (u, w)) \supset
        \exists t(R(v, t) \land R(w, t)))
```

We will investigate only the ones in red.

Modal logics vs. properties of accessibility relations

Κ		the class of all frames
K4	4	the class of transitive frames
КТ	Т	the class of reflexive frames
КВ	В	the class of symmetric frames
KD		the class of serial frames
KT4	<b>S4</b>	the class of reflexive and transitive frames
KT4B	<b>S</b> 5	the class of frames with an equivalence relation
KT5	S5	the class of frames with an equivalence relation
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